

$$1. \quad \underline{(i-j)m_k} + \underline{(j-k)m_i} + \underline{(k-i)m_j} = C, \quad (1)$$

在①中用 j, i 代替 i, j 得 $\underline{(j-i)m_k} + \underline{(i-k)m_j} + \underline{(k-j)m_i} = C \quad (2)$

①+②可得 $2C=0 \therefore C=0 \therefore$ ①化为

$$(i-j)m_k + (j-k)m_i + (k-i)m_j = 0 \quad (3)$$

在①中取 i, j, k 分别为 $1, 2, 3$ 可得

$$(1-2) \cdot \frac{a_1+a_2+a_3}{3} + (2-3)a_1 + (3-1) \cdot \frac{a_1+a_2}{2} = 0 \Rightarrow a_1+a_3=2a_2$$

假设 $a_1, a_2, \dots, a_n$ 为等差数列 ($n \geq 4$)

在①中取 i, j, k 为 $1, 2, n$ 可得

$$(1-2) \cdot \frac{a_1+a_2+\dots+a_n}{n} + (2-n)a_1 + (n-1) \cdot \frac{a_1+a_2}{2} = 0 \quad (4)$$

设 $a_i = a_1 + (a_2 - a_1) \cdot (i-1)$ ($i = 1, 2, \dots, n-1$) 代入④

可得 $a_n = a_1 + (a_2 - a_1)(n-1) \therefore$ 数列 $a_1, a_2, \dots$ 为等差数列。



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$$2. \because \text{第 } 4 \text{ 行与第 } 1 \text{ 列的比为 } \frac{3}{16} - \frac{1}{8} = \frac{1}{16} \therefore a_{4j} = \frac{j}{16} \quad (j=1, 2, \dots, n)$$

$$\therefore a_{44} = \frac{4}{16} = \frac{1}{4} \quad \text{设第 } 3 \text{ 行与第 } 4 \text{ 列的比为 } q. \text{ 则 } q^2 = \frac{a_{44}}{a_{24}} = \frac{1}{4} \therefore q = \frac{1}{2}$$

$$\therefore a_{ij} = \frac{j}{16} \cdot \left(\frac{1}{2}\right)^{i-4} = \frac{j}{2^i} \quad (i, j=1, 2, \dots, n)$$

$$\therefore a_{kk} = \frac{k}{2^k} \quad (k=1, 2, \dots, n)$$

$$\text{设 } S = \sum_{k=1}^n a_{kk} = \frac{1}{2} + \frac{2}{2^2} + \dots + \frac{n-1}{2^{n-1}} + \frac{n}{2^n} \quad \textcircled{1}$$

$$\frac{1}{2}S = \frac{1}{2^2} + \dots + \frac{n-2}{2^{n-1}} + \frac{n-1}{2^n} + \frac{n}{2^{n+1}} \quad \textcircled{2}$$

$$\textcircled{1} - \textcircled{2} \text{ 得 } \frac{S}{2} = \frac{1}{2} + \frac{1}{2^2} + \dots + \frac{1}{2^n} - \frac{n}{2^{n+1}}$$

$$\therefore S = 1 + \frac{1}{2} + \dots + \frac{1}{2^{n-1}} - \frac{n}{2^n}$$

$$= \frac{1 - \frac{1}{2^n}}{1 - \frac{1}{2}} - \frac{n}{2^n} = 2 - \frac{n+2}{2^n}$$



3. $n=3$ 时 $a_i=i (i=1,2,3)$ 则 $a_i+a_j=3,4,5$ ✓

$n=4$ 时 $a_i=i (i=1,2,3), a_4=5$, 则 $a_i+a_j=3,4,5,6,7,8$. ✓

$n \geq 5$ 时. 假设存在 $a_1, a_2, \dots, a_n$ 满足条件. 不妨假设 $a_1 < a_2 < \dots < a_n$

① 将 $a_i+a_j (1 \leq i < j \leq n)$ 由小到大排列为公差为 d 的等差数列 S . 且 $d > 0$.

则 S 的前两项为 a_1+a_2, a_1+a_3 . 最后两项为 $a_{n-2}+a_n, a_{n-1}+a_n$

$$\therefore a_3 - a_2 = (a_1+a_3) - (a_1+a_2) = d = (a_{n-1}+a_n) - (a_{n-2}+a_n) = a_{n-1} - a_{n-2} > 0$$

当 $n \geq 6$ 时 可得 $a_3 + a_{n-2} = a_2 + a_{n-1} \therefore S$ 中有两项相等 矛盾

当 $n=5$ 时 则 S 中 $C_5^2 = 10$ 项. 最后两项为 a_3+a_5, a_4+a_5 .

$$d = a_3 - a_2 = a_4 - a_3. \therefore a_1+a_4 - (a_1+a_3) = d \therefore a_1+a_4 \text{ 为 } S \text{ 中第 } 3 \text{ 项}$$

$$a_2+a_5 - (a_3+a_5) = d \therefore a_2+a_5 \text{ 为 } S \text{ 中第 } 8 \text{ 项}$$

$$\therefore a_3+a_4 - (a_2+a_4) = d = (a_2+a_4) - (a_2+a_3)$$

$\therefore a_2+a_3, a_2+a_4, a_3+a_4$ 是 S 中相邻的三项



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若 $a_1 + a_5$ 是 S 中第 4 次. 则

$$a_5 - a_4 = a_1 + a_5 - (a_1 + a_4) = d = a_3 - a_2 \therefore a_2 + a_5 = a_3 + a_4, \text{ 矛盾}$$

若 $a_1 + a_5$ 是 S 中第 7 次. 则

$$a_2 - a_1 = a_2 + a_5 - (a_1 + a_5) = d = a_4 - a_3 \therefore a_2 + a_3 = a_1 + a_4, \text{ 矛盾}$$

4. (1) $\Rightarrow$. $\because \frac{a_{n+1}}{a_n} \geq q > 1 \therefore a_{n+1} > a_n \therefore \{a_n\}$ 严格递增.

取 $r = q > 1$ 下证对任意 $x > 0$ 区间 (x, rx) 中

至多包含数列 $\{a_n\}$ 中的一项

假设 (x, rx) 中有 $\{a_n\}$ 中两项. 不妨假设

$$x < a_m < a_{m+1} < rx \therefore \frac{a_{m+1}}{a_m} < \frac{rx}{x} = r = q. \text{ 矛盾}$$

$\therefore \{a_n\}$ 是 S 数列

121 证.



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设 $b_n = 2^{n-1}$, ($n \in \mathbb{Z}_+$), 则 $\{b_n\}$ 是 L 数列 / 且 $\{b_n\}$ 是 S 数列
 设 $\begin{cases} a_{2i} = b_{2i-1} \\ a_{2i+1} = b_{2i} \end{cases}$ ($i \in \mathbb{Z}_+$), 则 $\{a_n\}$ 是 S 数列
 $\therefore \{a_n\}$ 不是齐次递推关系 $\therefore \{a_n\}$ 不是 L 数列

5 已知 a_1, a_2 , 且 $a_{n+2} + pa_{n+1} + qa_n = 0$, $n \in \mathbb{Z}_+$ = 二阶齐次递推关系

特征方程: 特征方程: $x^2 + px + q = 0$. 两个根 x_1, x_2 称为特征根

(1) 若 $x_1 \neq x_2$ 则 $a_n = C_1 x_1^n + C_2 x_2^n$. 其中 C_1, C_2 由
 $a_1 = C_1 x_1 + C_2 x_2$, $a_2 = C_1 x_1^2 + C_2 x_2^2$ 确定

定义 a_0 满足 $a_2 + pa_1 + qa_0 = 0$ ($q \neq 0$). $a_0 = C_1 + C_2$
 (2) 若 $x_1 = x_2$. 则 $a_n = (C_1 + C_2 n) x_1^n$. 其中 C_1, C_2 由
 $a_0 = C_1$, $a_1 = (C_1 + C_2) x_1$ 确定



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$$f_n = \frac{1}{\sqrt{5}} \left(\left(\frac{1+\sqrt{5}}{2} \right)^n - \left(\frac{1-\sqrt{5}}{2} \right)^n \right)$$

特征根为 $\frac{1+\sqrt{5}}{2}$, $\frac{1-\sqrt{5}}{2}$. 特征方程为 $x^2 - x - 1 = 0$

$$\therefore f_{n+2} - f_{n+1} - f_n = 0 \quad (n \geq 0) \quad f_0 = 0 \quad f_1 = 1$$

$\Rightarrow$ 所求为线性递推数列. 已证 a_0, a_1, a_2 且

$$a_{n+3} + p a_{n+2} + q a_{n+1} + r a_n = 0 \quad (n \geq 0)$$

特征方程 $x^3 + p x^2 + q x + r = 0$. $\Rightarrow$ 三个根 x_1, x_2, x_3 为特征根

(1) 若 x_1, x_2, x_3 两两不等. 则 $a_n = C_1 x_1^n + C_2 x_2^n + C_3 x_3^n$.

(2) 若 $x_1 = x_2 \neq x_3$. 则 $a_n = (C_1 + C_2 n) x_1^n + C_3 x_3^n$.

(3) 若 $x_1 = x_2 = x_3$. 则 $a_n = (C_1 + C_2 n + C_3 n^2) x_1^n$.

$$a_{n-2} a_{n+1} = 3 + a_n a_{n-1} \quad (1) \quad a_n + a_{n+2} = 3 + a_{n+1} a_n \quad (2)$$

$$(2) - (1) \quad a_n + a_{n+2} - a_{n-2} a_{n+1} = a_{n+1} a_n - a_n a_{n-1}$$



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$$a_{n-1}(a_n + a_{n+2}) = a_{n+1}(a_{n-2} + a_n) \quad (3)$$

$$\therefore a_n > 0$$

$$\therefore \text{由 (3) 得 } \frac{a_n + a_{n+2}}{a_{n+1}} = \frac{a_{n-2} + a_n}{a_{n-1}} \quad (4)$$

$$\sqrt[2]{b_n} = \frac{a_n + a_{n+2}}{a_{n+1}} \quad (n \in \mathbb{Z}_+)$$

$$\text{由 (4) 得 } b_n = b_{n-2} \quad \because a_4 = 5 \therefore b_1 = b_2 = 3 \quad \therefore b_n = 3 \quad (n \in \mathbb{Z}_+)$$

$$\therefore \frac{a_n + a_{n+2}}{a_{n+1}} = 3 \quad \therefore a_{n+2} - 3a_{n+1} + a_n = 0$$

$$\therefore a_n = \frac{5-2\sqrt{5}}{5} \left(\frac{3+\sqrt{5}}{2}\right)^n + \frac{5+2\sqrt{5}}{5} \left(\frac{3-\sqrt{5}}{2}\right)^n$$

$$6. \quad a_4 = 9, \quad b_1 = 3, \quad b_2 = 5 \quad \therefore b_{2k-1} = 3, \quad b_{2k} = 5 \quad (k \in \mathbb{Z}_+) \quad \text{定义 } b_n = \frac{a_n + a_{n+2}}{a_{n+1}}$$

$$\therefore 3 = b_{2k-1} = \frac{a_{2k-1} + a_{2k+1}}{a_{2k}} \Rightarrow a_{2k-1} + a_{2k+1} = 3a_{2k} \quad (1)$$

$$5 = b_{2k} = \frac{a_{2k} + a_{2k+2}}{a_{2k+1}} \Rightarrow a_{2k} + a_{2k+2} = 5a_{2k+1} \quad (2)$$

$$\text{由 (2) 得 } a_{2k} = \frac{1}{3}(a_{2k-1} + a_{2k+1}) \quad \text{由 (1) 得}$$



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$$\frac{1}{3}(a_{2k-1} + a_{2k+1}) + \frac{1}{3}(a_{2k+1} + a_{2k+3}) = 5a_{2k+1}$$

$$\therefore a_{2k+3} - 13a_{2k+1} + a_{2k-1} = 0 \quad a_{2k-1} = x_k$$

$$(8) \text{ 令 } a_{2k+2} - 13a_{2k} + a_{2k-2} = 0 \quad a_{2k} = y_k$$

$$\therefore \text{对 } \forall k \in \mathbb{Z}_+, a_{n+4} - 13a_{n+2} + a_n = 0.$$

~~由数字归纳法可知 a_n 为整数~~

7. 由第一个等式可得 $b_n = \frac{a_{n+1} - 7a_n + 3}{6}$ 代入第二个等式

$$\frac{a_{n+2} - 7a_{n+1} + 3}{6} = 8a_n + 7 \cdot \frac{a_{n+1} - 7a_n + 3}{6}$$

$$\therefore a_{n+2} - 14a_{n+1} + a_n + 6 = 0 \quad \text{待定 } \alpha \quad \text{则}$$

$$a_{n+2} - \alpha - 14(a_{n+1} - \alpha) + a_n - \alpha = 0 \quad \text{令 } 12\alpha = 6 \quad \alpha = \frac{1}{2}$$

$$\text{设 } A_n = a_n - \frac{1}{2} \quad \text{则 } A_{n+2} - 14A_{n+1} + A_n = 0$$



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$$\therefore A_n = \frac{1}{4}(7+4\sqrt{3})^n + \frac{1}{4}(7-4\sqrt{3})^n$$

$$\therefore a_n = \frac{1}{4}(7+4\sqrt{3})^n + \frac{1}{4}(7-4\sqrt{3})^n + \frac{1}{2}$$

$$= \left(\frac{(2+\sqrt{3})^n}{2} \right)^2 + \left(\frac{(2-\sqrt{3})^n}{2} \right)^2 + 2 \cdot \frac{(2+\sqrt{3})^n}{2} \cdot \frac{(2-\sqrt{3})^n}{2}$$

$$= \left(\frac{(2+\sqrt{3})^n}{2} + \frac{(2-\sqrt{3})^n}{2} \right)^2$$

$$\therefore \frac{(2+\sqrt{3})^n}{2} + \frac{(2-\sqrt{3})^n}{2} \in \mathbb{Z}_+ \therefore a_n \text{ 为完全平方数}$$

8 (1) $\because a_n > 0 \therefore a_{n+1} > \frac{7a_n}{2} > a_n \therefore \{a_n\}$ 严格递增

$$\therefore 2a_{n+1} - 7a_n = \sqrt{45a_n^2 - 36} \therefore 4a_{n+1}^2 + 49a_n^2 - 28a_na_{n+1} = 45a_n^2 - 36$$

$$\therefore a_{n+1}^2 + a_n^2 - 7a_na_{n+1} + 9 = 0 \quad ① \quad a_{n+2}^2 + a_{n+1}^2 - 7a_{n+1}a_{n+2} + 9 = 0 \quad ②$$

② - ① 得

$$(a_{n+2} - a_n)(a_{n+2} + a_n - 7a_{n+1}) = 0$$

$$\because a_{n+2} > 7a_{n+1} - a_n \text{ 由数学归纳法可得}$$



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a_n 为有理数 (2) 由①得 $a_{n+1}^2 + a_n^2 + 2a_n a_{n+1} = 9(a_n a_{n+1} - 1)$

$\therefore a_n a_{n+1} - 1 = \frac{(a_n + a_{n+1})^2}{3}$ ③

$\because a_n a_{n+1} - 1 \in \mathbb{Z}, \frac{a_n + a_{n+1}}{3} \in \mathbb{Q} \therefore \text{由③得 } \frac{a_n + a_{n+1}}{3} \in \mathbb{Z},$

$\therefore \text{由③得 } a_n a_{n+1} - 1 \text{ 为完全平方数}$

9. 设 $t = \lambda^{\frac{1}{3}} > 0$ 则 $t^3 = t + 1 \therefore t^3 = t^2 + 2t + 1$ ①

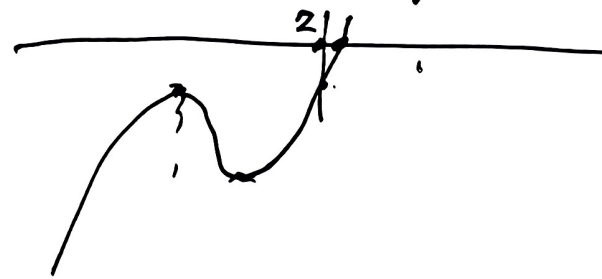
设 $f(x) = x^3 - x^2 - 2x - 1$ 则 $f'(x) = 3x^2 - 2x - 2$ 令 $f'(x) = 0$ 得到

解 $x_1 = \frac{1+\sqrt{17}}{3}, x_2 = \frac{1-\sqrt{17}}{3}$

$\therefore f(1) < 0, f(2) < 0 \therefore t > 2$

设 $f(x) = 0$ 的不同于 t 的 -2 个复根

为 $a, \bar{a}$ 由韦达定理得 $t + a + \bar{a} = 1$.



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$$|a| = |\bar{a}| = \frac{1}{t^{\frac{1}{2}}} < \frac{1}{2^{\frac{1}{2}}}$$

设 $u_n = t^n + a^n + \bar{a}^n$. ($n=0, 1, \dots$) $\frac{1}{2} \sqrt{2} \sqrt{2} \sqrt{2}$

$$u_{n+3} - u_{n+2} - 2u_{n+1} - u_n = 0 \quad (n=0, 1, \dots) \quad (3)$$

$$\therefore u_0 = 3, u_1 = t + a + \bar{a} = 1, u_2 = (t + a + \bar{a})^2 - 2(ta + t\bar{a} + a\bar{a})$$

$$= 1^2 - 2(-2) = 5 \quad (2) \text{ 及 } \bar{z} = 1/2 \text{ 的 } 1/3 \text{ 的 } u_n \in \mathbb{Z}_+ =$$

$$\begin{aligned} \overline{\lambda} M &= u_{450} \in \mathbb{Z}_+ \quad \text{p.} \quad |M - \lambda^{300}| = |u_{450} - t^{450}| \\ &= |a^{450} + \bar{a}^{450}| \leq |a|^{450} + |\bar{a}|^{450} < 2 \cdot \frac{1}{2^{\frac{450}{2}}} = \frac{1}{2^{225}} \\ &< \frac{1}{2^{200}} = \frac{1}{4^{100}}. \end{aligned}$$



$$10 \quad a_n + a_{n+1} = 2a_{n+2} a_{n+3} + 2024 \quad (1) \quad \therefore a_{n+1} + a_{n+2} = 2a_{n+3} a_{n+4} + 2024 \quad (2)$$

$$(2) - (1) \text{ 得 } \underline{a_{n+2} - a_n} = 2a_{n+3} (\underline{a_{n+4} - a_{n+1}}) = \dots = 2^k a_{n+3} a_{n+5} \dots a_{n+2k+1} (\underline{a_{n+2k+2} - a_{n+1}}) \quad (3)$$

其中 k 为任意正整数
对于任意固定正整数 n

(由 (3) 得 $2^k \mid a_{n+2} - a_n$, k 为任意正整数
 $\therefore a_{n+2} - a_n = 0$. 故 $a_1 = a_3 = \dots = a_{2t-1} = \dots \triangleq a$, $a_2 = a_4 = \dots = a_{2t} = \dots \triangleq b$ (由 (1) 得 $a+b = 2ab + 2024$)

$$2a+2b = 2a \cdot 2b + 4048 \quad \therefore (2a-1)(2b-1) = -4047 = -3 \times 19 \times 71.$$

$$2a-1 = \pm 1, \pm 3, \pm 19, \pm 71, \pm 57, \pm 213, \pm 1349, \pm 4047$$

$\therefore (a_1, a_2)$ 有 16 种取值

$$11. \quad a_{n+1} a_{n+2} + a_{n+2} = a_n + 2024 \quad (1) \quad a_{n+2} a_{n+3} + a_{n+3} = a_{n+1} + 2024 \quad (2)$$

$$(2) - (1) \text{ 得 } \underline{a_{n+2} a_{n+3}} - \underline{a_{n+1} a_{n+2}} + \underline{a_{n+3}} - \underline{a_{n+2}} = \underline{a_{n+1}} - \underline{a_n}$$



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$$a_{n+2}(a_{n+3}-a_{n+1}) + a_{n+3}-a_{n+1} = a_{n+2}-a_n$$

$$\therefore (a_{n+2}+1)(a_{n+3}-a_{n+1}) = a_{n+2}-a_n \quad (2) \quad \text{其中 } a_{n+2}+1 \neq 0.$$

$$\text{若 } a_1=a_3, \text{ 由 (2) 可得 } a_2=a_4, a_3=a_5, \dots$$

$$\therefore a_1=a_3=\dots=a_{2b-1}=\dots=a, a_2=a_4=\dots=a_{2b}=\dots=b.$$

$$(t \in \mathbb{Z}) \text{ 代入 (1) 可得 } ab = 2024 = 2^3 \cdot 11 \cdot 23$$

$$\therefore \{a_n\} \text{ 个数为 } 2 \cdot (3+1)(1+1)(1+1) - 2 = 30.$$

$$\text{若 } a_1 \neq a_3 \text{ 由 (2) 可得 } a_2 \neq a_4, a_3 \neq a_5, \dots \text{ 即 2 个情况}$$

$$\text{正整数 } n \text{ 总有 } a_n \neq a_{n+2} \therefore a_{n+3}-a_{n+1} \neq 0, a_{n+2}-a_n \neq 0 (n \in \mathbb{Z}_+)$$

$$\text{由 (2) 可得 } |a_{n+2}-a_n| = |a_{n+2}+1| |a_{n+3}-a_{n+1}| \geq |a_{n+3}-a_{n+1}|$$

$$\therefore \text{正整数 } n, \{ |a_{n+2}-a_n| \} \text{ 单调递增} \therefore \text{有 } k \text{ 个正整数}$$



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当 $n \geq N$ 时 $|a_{n+2} - a_n| = C \in \mathbb{Z}_+$, C 为常数

当时 $|a_{n+2} - a_n| = 1 \quad \therefore a_{n+2} \in \{0, -2\}$

对于 $n \geq N+2$ 有 $a_n, a_{n+1}, a_{n+2} \in \{0, -2\}$.

不满足①.

14

