

当 $n \geq N$ 时 $|a_{n+2} - a_n| = C \in \mathbb{Z}_+$. C 为常数

这时 $|a_{n+2} + 1| = 1 \quad \therefore a_{n+2} \in \{0, -2\}$

对于 $n \geq N+2$ 有 $a_n, a_{n+1}, a_{n+2} \in \{0, -2\}$.

不满足①.

(2. $f(1)=f(2)=1, f(3)=2$.

(1) 若 $a_2=2$. 设 $b_i = a_{i+1} - 1$ ($i=1, 2, \dots, n-1$) 有 $b_1, b_2, \dots, b_{n-1} \in \{1, 2, \dots, n-1\}$
 $a_1 - 1$ 排列. $b_1=1$. 且 $|b_i - b_{i-1}| = |a_{i+1} - a_i| \leq 2 \therefore$ 满足条件有
 $f(n-1) \uparrow$

(2) 若 $a_2=3$ (i) 当 $a_3=2$ 时 $a_4=4 \dots$ 设 $c_i = a_{i+3} - 3$ ($i=1, 2, \dots, n-3$)
 有 $c_1, c_2, \dots, c_{n-3} \in \{1, 2, \dots, n-3\}$ 且 $|c_i - c_{i-1}| = |a_{i+3} - a_{i+2}| \leq 2 \therefore$ 满足条件有 $f(n-3) \uparrow$



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(ii) $a_3=5, \dots, \underline{a_k=2k-1}$ a_{k+1} 为偶数

若 $a_{k+1}=2k$ ($k \geq 2$) 若 $a_{k+2}=2k+1$ 或 $2k+2$. k 在排列中不出现 2.

出现 2. 矛盾. $\therefore a_{k+2}=2k-2, \dots, a_{2k}=2$. 此时

$n=2k$ 此时在排列中只有 1 个

若 $a_{k+1}=2k-2$ ($k \geq 3$). 若 $a_{k+2}=2k$, k 在排列中不出现 2.

矛盾. $a_{k+2}=2k-4, \dots, a_{2k-1}=2$. 此时 $n=2k-1$ 此时

在排列中只有 1 个. $\therefore f(n)=f(n-1)+f(n-3)+1$. ($n \geq 4$)

$\therefore f(n) \equiv 1, 1, 2, 0, 2, 1, 2, 1, 3, 2, 0, 0, 3, 0, 1, 1, 2 \pmod{4}$

$f(n)$ 模 4 在 18 以内为 14. $\therefore 2024 \equiv 8 \pmod{14}$,

$\therefore f(2024) \equiv f(8) \equiv 1 \pmod{4}$



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$$\therefore a_n(b_n + b_{n+2}) = b_{n+1}$$

$$\therefore a_n(b_{n+2}^2 - b_n^2) = a_n(b_n + b_{n+2})(b_{n+2} - b_n) = \underline{b_{n+1}(b_{n+2} - b_n)}$$

$$= -\frac{1}{2}((b_{n+1} - b_{n+2})^2 - (b_{n+1}^2 + b_{n+2}^2)) + \frac{1}{2}((b_n - b_{n+1})^2 - (b_n^2 + b_{n+1}^2))$$

$$(2a_n - 1)(b_{n+2}^2 - b_n^2) = (b_n - b_{n+1})^2 - (b_{n+1} - b_{n+2})^2 \quad ①$$

$$\therefore 2a_n - 1 \geq 2a_0 - 1 > 2 \times \frac{1}{2} - 1 = 0$$

$$\therefore \text{由①得 } b_{n+2}^2 - b_n^2 = \frac{(b_n - b_{n+1})^2 - (b_{n+1} - b_{n+2})^2}{2a_n - 1} \quad ②$$

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在②中令 $n = 0, 1, \dots, n-2$ 并求和

$$b_{n+2}^2 - b_n^2 - b_0^2 + b_1^2 = \frac{(b_0 - b_1)^2}{2a_0 - 1} - \frac{(b_1 - b_2)^2}{2a_0 - 1} + \frac{(b_1 - b_2)^2}{2a_1 - 1} - \frac{(b_2 - b_3)^2}{2a_1 - 1} +$$

$$\dots + \frac{(b_{n-2} - b_{n-1})^2}{2a_{n-2} - 1} - \frac{(b_{n-1} - b_n)^2}{2a_{n-2} - 1}$$

$$= \frac{(b_0 - b_1)^2}{2a_0 - 1} - (b_1 - b_2)^2 \left(\frac{1}{2a_0 - 1} - \frac{1}{2a_1 - 1} \right) - \dots - (b_{n-2} - b_{n-1})^2 \left(\frac{1}{2a_{n-3} - 1} - \frac{1}{2a_{n-2} - 1} \right) - \frac{(b_{n-1} - b_n)^2}{2a_{n-2} - 1}$$



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$$\leq \frac{(b_0 - b_1)^2}{2a_0 - 1}$$

$$\therefore b_n^2 \leq b_n^2 + b_{n-1}^2 \leq b_0^2 + b_1^2 + \frac{(b_0 - b_1)^2}{2a_0 - 1} \triangleq |M|^2$$

$$\therefore |b_n| \leq |M|$$

14. 分析. 假设 $\{x_n\}$ 是常数列. 设 $x_n = a > 0$ 且 a 是方程

$$x^2 = \alpha x - a \quad \therefore a^2 = \alpha(a-1) \quad \therefore a = \alpha - 1 > 0 \quad \therefore \alpha > 1$$

若 $\alpha > 1$, 则 $x_n = \alpha - 1 > 0$ 且 $a^2 = \alpha(a-1) - (a-1)$ $\square$

$= \alpha - 1$, 是方程 $\checkmark$

若 $\alpha \leq 1$ 且存在常数列 $\{x_n\}$.

$\therefore \alpha x_{n+1} - x_n > 0 \quad \therefore x_{n+1} \geq \alpha x_{n+1} > x_n \quad \therefore \{x_n\}$ 严格递增

$$\therefore x_{n+2}^2 = \alpha x_{n+1} - x_n \stackrel{\textcircled{1}}{\leq} x_{n+1} - x_n < x_{n+1}$$

$$\therefore x_{n+1}^2 \leq x_{n+2}^2 < x_{n+1} \quad \therefore x_{n+1} < 1$$



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在①中①1, 2, ..., n 代 x_n . 并求和得1/4

$$1) x_3^2 < x_3^2 + x_4^2 + \dots + x_{n+2}^2 \leq x_{n+1} - x_1 < x_{n+1} \quad (2)$$

当 $n \rightarrow +\infty$ ②与 $x_{n+1} < 1$ 矛盾

15 $x_{n+2} - (k+1)x_{n+1} + x_n = 0$ 特征方程为 $\lambda^2 - (k+1)\lambda + 1 = 0$ ①

若 $\lambda = 1$ 则 $1 - k - 1 + 1 = 0 \Rightarrow k = 1$ 矛盾 $\therefore \lambda \neq 1$

设①两根为 $\alpha, \frac{1}{\alpha}$ 且 α 为实数设 $\alpha > 1$.

由韦达定理得 $\alpha + \frac{1}{\alpha} = k+1$

$$\alpha = \frac{k+1 + \sqrt{(k+1)^2 - 4}}{2}$$

设 $x_n = C_1 \alpha^n + C_2 \alpha^{-n}$

$\therefore x_0 = 1 = C_1 + C_2$

$\therefore C_1 = \frac{\alpha}{\alpha-1} \quad C_2 = -\frac{1}{\alpha-1}$

$\alpha + \frac{1}{\alpha} + 1 = k+2 = x_1 = C_1 \alpha + C_2 \cdot \frac{1}{\alpha}$

$\therefore x_n = \frac{\alpha^{n+1}}{\alpha-1} - \frac{\alpha^{-n}}{\alpha-1} \quad (2)$

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$$\alpha = \frac{2k+2 + 2\sqrt{(k-1)(k+3)}}{4} = \left(\frac{\sqrt{k-1} + \sqrt{k+3}}{2} \right)^2 = \beta^2.$$

$$\text{例 2) 证明 } x_n = \frac{\beta^{2n+2} - \beta^{-2n}}{\beta^2 - 1} = \frac{\beta^{2n+1} - \beta^{-(2n+1)}}{\beta - \beta^{-1}} \quad (3)$$

$$\text{设 } k = m^2 - 3 \quad (m \in \mathbb{Z}_+, m \geq 3) \quad \text{则}$$

$$\beta = \frac{\sqrt{k-1} + \sqrt{k+3}}{2} = \frac{\sqrt{m^2-4} + m}{2}$$

$$= \frac{2\sqrt{(m-2)(m+2)} + 2m}{2} = \left(\frac{\sqrt{m-2} + \sqrt{m+2}}{2} \right)^2 = \gamma^2$$

$$\text{例 3) 证明 } x_n = \frac{\gamma^{2(2n+1)} - \gamma^{2(2n+1)^{-2}}}{\gamma^2 - \gamma^{-2}} = \frac{\gamma^{2n+1} - \gamma^{-(2n+1)}}{\gamma - \gamma^{-1}} \cdot \frac{\gamma^{2n+1} + \gamma^{-(2n+1)}}{\gamma + \gamma^{-1}}$$

$$= y_n z_n \quad \therefore y_n = \frac{\gamma}{\gamma - \gamma^{-1}} \cdot \beta^n = \frac{\gamma^{-1}}{\gamma - \gamma^{-1}} \cdot \beta^{-n}.$$

$$\begin{aligned} y_{n+2} - m y_{n+1} + y_n &= 0 \quad (4) & y_0 &= 1, & y_1 &= m+1. \\ z_{n+2} - m z_{n+1} + z_n &= 0 \quad (5) & z_0 &= 1, & z_1 &= m-1 \geq 2 \end{aligned}$$



④⑤ 已知 $y_n, z_n \in \mathbb{Z}_+$ 且当 $n \in \mathbb{Z}_+$ 时 $y_n, z_n \geq 2$

$\therefore x_n$ 是合数

$x_0 = 1$ 不是素数

$$(6) \quad x_{n+1} = \frac{x_n^2}{x_n^2 - x_n + 1} = 1 + \frac{x_n - 1}{x_n^2 - x_n + 1} \quad \therefore x_{n+1} - 1 = \frac{x_n - 1}{x_n^2 - x_n + 1} \quad (1)$$

$$\because x_n^2 - x_n + 1 > 0 \quad \text{且} \quad x_n - 1 \leq x_{n+1} - 1 \quad \text{同} \quad \therefore x_1 = \frac{4}{3} > 1$$

$\therefore$ 对任意 $n \in \mathbb{Z}_+$ 均有 $x_n > 1$

$$\text{由 (1) 得} \quad \frac{1}{x_{n+1} - 1} = \frac{x_n(x_n - 1) + 1}{x_n - 1} = x_n + \frac{1}{x_n - 1}$$

$$\therefore x_n = \frac{1}{x_{n+1} - 1} - \frac{1}{x_n - 1} \quad (2)$$

在 (2) 中 n 取 $1, 2, \dots, n$ 并求和

$$\begin{aligned} \text{并求和得} \quad \sum_{k=1}^n x_k &= \frac{1}{x_{n+1} - 1} - \frac{1}{x_1 - 1} = \frac{1}{x_{n+1} - 1} - \frac{1}{\frac{4}{3} - 1} = \frac{1}{x_{n+1} - 1} - 3 \\ &= \frac{x_n(x_n - 1) + 1}{x_n - 1} - 3 = \frac{x_n^2 - 4x_n + 4}{x_n - 1} = \frac{(x_n - 2)^2}{x_n - 1} \quad (3) \end{aligned}$$

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$$x_n \in \mathbb{Q}_+ \quad \therefore x_n - 1 \in \mathbb{Q}_+ \quad \text{设 } x_n - 1 = \frac{a}{b} \quad a, b \in \mathbb{Z}_+, (a, b) = 1$$

$$\text{式 } \lambda(3) \text{ 中 } \sum_{k=1}^n x_k = \frac{(\frac{a}{b} - 1)^2}{\frac{a}{b}} = \frac{(a-b)^2}{ab} \quad (4)$$

$$\because (a-b, a) = (a-b, b) = 1 \quad \therefore (4) \text{ 在是最简分数, 分子 } (a-b)^2$$

$$17.(1) \text{ 设 } s = \min \{a_{n+1} - a_n \mid n \in \mathbb{Z}_+\} \in \mathbb{Z}_+ \quad \text{对 } k \in \mathbb{Z}_+ \text{ 设 } 2^k > s$$

$$\text{设 } a_{m+1} - a_m = s \quad \text{且 } m \in \mathbb{Z}_+ \quad \text{且}$$

$$2(a_{2^k(m+1)} - a_{2^k m}) = a_{2^k(m+1)} - a_{2^k m} = \underbrace{a_{2^k(m+1)} - a_{2^k(m+1)-1}}_{\geq s} + \dots + \underbrace{a_{2^k m+1} - a_{2^k m}}_{\geq s}$$

$$\begin{aligned} & \vdots \\ & \parallel \\ & 2^k(a_{m+1} - a_m) \\ & \parallel \\ & 2^k s \end{aligned}$$

$$\geq s + \dots + s$$

$$\geq 2^k s$$

所以有

$$\text{即 } a_{2^k(m+1)} - a_{2^k(m+1)-1} = \dots = a_{2^k m+1} - a_{2^k m} = s$$



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$T = \{a_{2^k m+i} \mid i=0, 1, \dots, 2^k\}$ 为等差数列, 公差为 S

设 $T' = \{a_{2^k m+i} \mid i=0, 1, \dots, p-1\}$ 为等差数列, 公差为 S

假设存在 $0 \leq i < j \leq p-1$ 使得 $a_{2^k m+j} \equiv a_{2^k m+i} \pmod{p}$

$$k. \quad (j-i)S \equiv 0 \pmod{p} \quad \because 0 < j-i < p \quad \therefore S \equiv 0 \pmod{p}$$

$$\therefore 0 < S \leq a_2 - a_1 = 2a_1 - a_1 = a_1 < p \quad \text{矛盾}$$

$\therefore T'$ 是模 p 完全等差数列, $\therefore T'$ 中每一项是 p 的倍数

(2) 对于任意 $n \in \mathbb{Z}_+$ 存在唯一正整数 k_n , 使得 $2^{k_n} \leq n < 2^{k_n+1}$

$$\therefore 2^{k_{2n}} \leq 2n < 2^{k_{2n}+1} \quad \text{且} \quad 2^{k_n+1} \leq 2n < 2^{k_n+2} \quad \therefore k_{2n} = k_n + 1$$

且 $k_n \leq k_{n+1}$ 设 $a_n = 2^{k_n} + nP \quad \therefore$ 数列 $\{a_n\}$ 严格

递增 $\therefore a_{2n} = 2^{k_{2n}} + 2nP = 2^{k_n+1} + 2nP = 2(2^{k_n} + nP) = 2a_n$



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$\{a_n\}$ 是递增数列 $\therefore p \nmid a_n, (n \in \mathbb{Z}_+)$

18. 设 $r_i = \sqrt{\sigma(i)} + \sqrt{\dots + \sqrt{\sigma(n)}} > 0 (i=1, 2, \dots, n)$ $\because r_1 \in \mathbb{Q}_+$

$\therefore r_2 = r_1^2 - \sigma(1) \in \mathbb{Q}_+, \dots, r_n = r_{n-1}^2 - \sigma(n-1) \in \mathbb{Q}_+$

$\therefore r_n = \sqrt{\sigma(n)} \in \mathbb{Z}_+, r_{n-1} = \sqrt{\sigma(n-1) + r_n} \in \mathbb{Z}_+, \dots, r_1 = \sqrt{\sigma(1) + r_2} \in \mathbb{Z}_+$

下面证明 $\forall i=1, 2, \dots, n$ 均有 $r_i < \sqrt{n} + 1$ ①

$i=n$ 时 $r_n = \sqrt{\sigma(n)} \leq \sqrt{n} < \sqrt{n} + 1$ ✓

假设 $r_{i+1} < \sqrt{n} + 1$ 且 $r_i = \sqrt{\sigma(i) + r_{i+1}} < \sqrt{n + \sqrt{n} + 1} < \sqrt{n} + 1$ ✓

$\therefore$ ① 成立

对于任意正整数 n , 存在 $l^2 \leq n < (l+1)^2$

设 $\sigma(i) = l^2$

$i \in \{1, 2, \dots, n\}$

若 $i < n$ 且

$l < r_i = \sqrt{\sigma(i) + r_{i+1}} = \sqrt{l^2 + r_{i+1}} < \sqrt{n} + 1 < l + 2$



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$$\therefore r_i = l+1 \quad \therefore l^2 + r_{i+1} = l^2 + 2l + 1$$

$$\therefore 2l+1 = r_{i+1} < \sqrt{n+1} < l+2 \Rightarrow l < 1 \text{ 矛盾} \therefore i=1$$

$$\text{即 } \sigma(n) = l^2$$

$$\text{若 } l \geq 2, \quad \forall j \in \{1, 2, \dots, n-1\} \quad \text{使 } \sigma(j) = l^2 - 1$$

$$\therefore l < r_j = \sqrt{\sigma(j) + r_{j+1}} = \sqrt{l^2 - 1 + r_{j+1}} < \sqrt{n+1} < l+2$$

$$\therefore r_j = l+1 \quad \therefore l^2 - 1 + r_{j+1} = l^2 + 2l + 1 \quad r_{j+1} > 1$$

$$\therefore 2l+2 = r_{j+1} < \sqrt{n+1} < l+2 \Rightarrow l < 0 \text{ 矛盾}$$

$$\therefore l=1 \quad 1 \leq n \leq 4 \quad n=1, \sqrt{1} \checkmark$$

$$n=2, \sqrt{1+\sqrt{2}}, \sqrt{2+\sqrt{1}} \quad \times$$

$$n=3, \sqrt{2+\sqrt{3+\sqrt{1}}} \quad \checkmark$$

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19. 不妨假设 $x \geq y \geq z$ 则 $x^2 - 2(y+z)x + (y-z)^2 = 0$

$$x = \frac{2(y+z) \pm \sqrt{4(y+z)^2 - 4(y-z)^2}}{2} = y+z \pm 2\sqrt{yz} = (\sqrt{y} \pm \sqrt{z})^2$$

$$\therefore \sqrt{x} = \sqrt{y} \pm \sqrt{z} \quad \therefore \sqrt{x} = \sqrt{y} + \sqrt{z}$$

$$\therefore \text{原式} (=) \frac{(\sqrt{y} + \sqrt{z})^2 + y + z}{3} \geq \sqrt[3]{2(\sqrt{y} + \sqrt{z})^2 yz} \quad (1)$$

若 $y=0$ 则 (1) 成立.

若 $y>0$ 设 $t = \frac{z}{y}$ 则 (1) 化为 $\frac{2(t+1) + 2\sqrt{t}}{3} \geq \sqrt[3]{2t(t+1)^2} \quad (2)$

$$(2) \text{左} = \frac{2t + t+1 + \sqrt{t+1}}{3} \geq \sqrt[3]{2t \cdot (t+1)(t+1)} = (2) \text{右}$$

20 (1) 若 m 为偶数 设 $m=2k$ $(k \in \mathbb{Z}^+)$ 不妨设 $a_1 \leq a_2 \leq \dots \leq a_k$

$$\text{设 } A_1 = \frac{\sum_{i=1}^k a_i}{k}$$

$$A_2 = \frac{\sum_{i=k+1}^{2k} a_i}{k}$$

$$\therefore A_1 \leq 10 \leq A_2$$



$$\sum_{i=1}^{2k} A_i = 10 \quad \text{且} \quad A_2 = 10 \quad \therefore a_i = 10 \quad (i=1, 2, \dots, 2k) \text{ 矛盾}$$

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$$\therefore A_1 < 10 \quad \text{且} \quad A_2 > 10$$

$$\text{设 } A_1 = 10 - x, \quad x > 0 \quad \text{且} \quad A_2 = 10 + x.$$

$$\text{若 } x < 1 \quad \text{且} \quad A_1 > 9 \quad A_2 < 11$$

$$\therefore a_k \geq 1 \quad \therefore a_{2k} \geq \dots \geq a_{k+1} \geq a_k \geq 1 \quad \therefore A_2 \geq 11 \text{ 矛盾}$$

$$\therefore x \geq 1 \quad \therefore \left(\frac{1}{2k} \sum_{i=1}^{2k} a_i \right)^{\frac{1}{2k}} = \sqrt{\left(\frac{1}{k} \sum_{i=1}^k a_i \right)^{\frac{1}{k}} \left(\frac{1}{k} \sum_{i=k+1}^{2k} a_i \right)^{\frac{1}{k}}} \\ \leq \sqrt{A_1 A_2} = \sqrt{(10-x)(10+x)} = \sqrt{100 - x^2} \leq \sqrt{100 - 1} = 3\sqrt{11}.$$

均值不等式

$$(2) \text{ 若 } m \text{ 为奇数,} \quad \sum_{i=1}^m a_i + \sum_{i=1}^m a_i = 10 \cdot 2m. \quad \therefore \text{ 平均 } 2m. \quad \text{由 (1)}$$

$$\text{所以 } \left(\frac{1}{m} \sum_{i=1}^m a_i \right)^{\frac{1}{m}} = \left(\frac{1}{m} \sum_{i=1}^m a_i \left(\frac{1}{m} \right)^{\frac{1}{m}} \right)^{\frac{1}{2m}} \leq 3\sqrt{11}$$

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$$21 \quad \frac{a + \sqrt{ab} + \sqrt[3]{abc} + \sqrt[4]{abcd}}{4} \leq \sqrt[4]{a \cdot \frac{a+b}{2} \cdot \frac{a+b+c}{3} \cdot \frac{a+b+c+d}{4}} \quad ①$$

$$① \Leftrightarrow \frac{1}{4} \left(\sqrt[4]{\frac{a}{a} \cdot \frac{2a}{a+b} \cdot \frac{3a}{a+b+c} \cdot \frac{4a}{a+b+c+d}} + \sqrt[4]{\frac{a}{a} \cdot \frac{2a}{a+b} \cdot \frac{3b}{a+b+c} \cdot \frac{4b}{a+b+c+d}} \right. \\ \left. + \sqrt[4]{\frac{a}{a} \cdot \frac{2b}{a+b} \cdot \frac{3\sqrt[3]{abc}}{a+b+c} \cdot \frac{4c}{a+b+c+d}} + \sqrt[4]{\frac{a}{a} \cdot \frac{2b}{a+b} \cdot \frac{3c}{a+b+c} \cdot \frac{4d}{a+b+c+d}} \right) \leq 1. \quad ②$$

② 均值不等式

$$② \text{ 左} \leq \frac{1}{4} \cdot \frac{1}{4} \left(\frac{a}{a} + \frac{2a}{a+b} + \frac{3a}{a+b+c} + \frac{4a}{a+b+c+d} + \frac{a}{a} + \frac{2a}{a+b} + \frac{3b}{a+b+c} + \frac{4b}{a+b+c+d} \right. \\ \left. + \frac{a}{a} + \frac{2b}{a+b} + \frac{3\sqrt[3]{abc}}{a+b+c} + \frac{4c}{a+b+c+d} + \frac{a}{a} + \frac{2b}{a+b} + \frac{3c}{a+b+c} + \frac{4d}{a+b+c+d} \right) \\ = \frac{1}{16} (4 + 4 + 3 + 4 + \frac{3\sqrt[3]{abc}}{a+b+c}) \leq \frac{1}{16} (15 + 1) = 1$$

均值不等式

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