

44 已知假设 $x > 0, y > 0, z = -t < 0$ $\therefore$ ~~已知~~ $f(x, y, t) = |x+y-t|$

在 $\frac{1}{2}$ 处 $f(u) = u + \frac{1}{u}$ ($u \geq 1$) $\therefore f(u)$ 单调递增. $\downarrow$

$$f(x) + f(y) = f(t) \quad \text{且} \quad f(t) = (x+y) + \frac{x+y}{xy} \geq (x+y) + \frac{1}{x+y} = f(x+y)$$

$$\therefore t > x+y, \therefore S(x, y, t) = t - (x+y), \quad (x+y)^2 > xy \quad \checkmark$$

3) 证. (1) 证 x, y, t 及 x', y', t' 若 $y' < y$ 且 $t - (x+y) < t' - (x'+y')$
 $f(x) + f(y) = f(t)$ ① $f(x) + f(y') = f(t')$ ② $S(x, y, t) \quad S(x', y', t')$

证 $\because f(y') < f(y)$ 由 ② ① 得 $f(t') < f(t) \Rightarrow \underline{t' < t}$

$$\because t > x+y > 1 \quad t' > x+y' > 1 \quad 1 < y't' < yt$$

$$1 > \frac{1}{y't'} > \frac{1}{yt} \quad \therefore 1 - \frac{1}{y't'} < 1 - \frac{1}{yt}$$

$$\text{由 ①, ②} \quad f(t) - f(y) = f(x) = f(t') - f(y')$$

$$(t-y)\left(1 - \frac{1}{ty}\right) = t-y + \frac{y-t}{ty}$$

$$t-y' + \frac{y'-t'}{ty'} = (t-y')\left(1 - \frac{1}{ty'}\right)$$



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$$\therefore t-y < t'-y' \Rightarrow S(x, y, t) < S(x, y', t')$$

① ② ③ ④ $S(x, y, t) \leq S(x, 1, t_1) \leq S(1, 1, t_2) = 2 + \sqrt{3} - 1 - 1 = \sqrt{3}$

⑤ x, y 在边界上

其中在边界上 $f(t_2) = f(1) + f(1)$ 即 $t_2 + \frac{1}{t_2} = 2 + 2 \therefore t_2 = 2 + \sqrt{3}$

② $x=1, y=1, z=-(2+\sqrt{3})$ 时 $(x+y+z)_{\max} = \sqrt{3}$.

45 $\sum \frac{x^2 y^2}{1-z} + 3xyz \leq k$ ①

① ② ③ $x=y=z=\frac{1}{3}$ 时 $k \geq \frac{1}{6} \therefore k = \frac{1}{6}$ 下证

$\sum \frac{x^2 y^2}{1-z} + 3xyz \leq \frac{1}{6}$ ②

② 左 = $xyz \left(\sum \frac{xy}{z(1-z)} + 3 \right) = xyz \cdot \sum \left(\frac{xy}{z(x+y)} + 1 \right)$

= $xyz \cdot \sum \frac{xy+yz+zx}{z(x+y)} = (xy+yz+zx) \cdot \sum \frac{xy}{x+y}$ ③

$\therefore (x+y+z)^2 \geq 3(xy+yz+zx) \therefore xy+yz+zx \leq \frac{1}{3}$ ④

④ ⑤ ⑥ ⑦ ⑧ $\sum \frac{xy}{x+y} \leq \frac{1}{4} \sum (x+y) = \frac{1}{2} \sum x = \frac{1}{4}$ ⑤. ④ ⑤ ⑥ ⑦ ⑧ $\leq \frac{1}{6}$

$\therefore R_{\min} = \frac{1}{6}$



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$$1. \quad 79|a+77b \Rightarrow 79|a-2b \Rightarrow 79|39a-78b \Rightarrow 79|39a+b$$

$$77|a+79b \Rightarrow 77|a+2b \Rightarrow 77|39a+78b \Rightarrow 77|39a+b \quad \because (77, 79)=1$$

$$\therefore 77, 79|39a+b \quad \text{设 } 39a+b=77 \cdot 79 \cdot k \quad (k \in \mathbb{Z}_+)$$

$$\therefore 39(a+b)=77 \cdot 79k+38b=(78^2-1)k+38b=(78^2-39)k+38(k+b).$$

$$\therefore 39|38(k+b) \quad \because (38, 39)=1 \quad \therefore 39|k+b$$

$$\therefore 39|a+b=(78^2-39)k+38(k+b) \geq (78^2-39) \cdot 1 + 38 \cdot 39$$

即 $k=1, k+b=39$ 中 $b=38$

$$\therefore a+b \geq 193. \quad \text{即 } a=155, b=38 \therefore (a+b)_{\min}=193$$

$$2. \quad \because p|n^2-1=(n-1)(n^2+n+1) \therefore \text{要 } p|n-1 \text{ 要 } p|n^2+n+1$$

$$\text{若 } p|n-1 \text{ 则 } n|p-1 \text{ 且 } p-1 \geq n > n-1 \geq p. \text{ 矛盾}$$

$$\therefore p|n^2+n+1 \quad \text{设 } n^2+n+1=tp \quad (t \in \mathbb{Z}_+), \quad p-1=ln, \quad l \in \mathbb{Z}_+$$

$$k.) \quad \underline{n^2+n+1} = \underline{t(l n + 1)} \quad \therefore 1 \equiv t \pmod{n} \quad \text{设 } t=mn+1. \quad (m \text{ 为非负整数})$$



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$$\text{式①} \text{ 中 } n^2+n+1 = (m+1)(l+1) \text{ ②}$$

若 $m \in \mathbb{Z}_+$ ②右 $\geq (n+1)(n+1) > \text{②左}$ 矛盾 $\therefore m=0, t=1$.

$$p = n^2+n+1 \quad \therefore 4p-3 = 4n^2+4n+1 = (2n+1)^2 \text{ 是奇数平方数.}$$

3. 若 $d=m-1$, 则对 $k \in \{m, m+1, \dots, n\}$ 总有 $k \nmid d$

$$\because n > m(m-2) \geq m > m-1 \quad \therefore d \nmid n! \quad \checkmark$$

故 下面证明 $d \leq m-1$ ①.

若 $d=1$, 则① $\checkmark$

若 $d \geq 2$ 设 $d = p_1 p_2 \dots p_t$ 其中 $p_1, p_2, \dots, p_t$ 为素数 由 $d \mid n!$ 知

$p_i \leq n \quad (i=1, 2, \dots, t)$ 若存在某个 $p_i \in \{m, m+1, \dots, n\}$

取 $k=p_i$ 则 $k \mid d$ 矛盾 $\therefore p_i \leq m-1 \quad (i=1, 2, \dots, t)$

特别地 $p_1 \leq m-1$, 假设存在 $j \in \{1, 2, \dots, t-1\}$ 使

$$p_1 p_2 \dots p_j \leq m-1, \quad p_1 p_2 \dots p_{j+1} \geq m \quad \text{由 } n \geq m(m-2)+1 = (m-1)^2$$



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$$k) \underline{p_1 p_2 \cdots p_j p_{j+1}} \leq (m-1) \cdot (m-1) \leq n.$$

若 $p_1 p_2 \cdots p_{j+1} \in \{m, m+1, \dots, n\}$. 取 $k = p_1 p_2 \cdots p_{j+1}$. 且 $k|d$, 矛盾
 $\therefore p_1 p_2 \cdots p_{j+1} \leq m-1$, 矛盾. 因此 $p_1 p_2 \cdots p_t = d \leq m-1$

$$\therefore d_{\max} = m-1$$

4. 若 $N = p^2 \cdot q$ 其中 p 为素数, $q \in \mathbb{Z}_+$ 则取 $x = pq, y = p^2 q - pq \in \mathbb{Z}_+$

5

$$\therefore x+y=N. \quad N|xy = pq \cdot (p^2 q - pq) = p^2 q(pq - q) \quad \checkmark$$

若 $N = p_1 p_2 \cdots p_t$ 其中 $p_1, p_2, \dots, p_t$ 是互不相同的素数 假设存在

正整数 x, y 使得 $x+y=N$. $N|xy$ 设 $xy = kN, k \in \mathbb{Z}_+$

$$\therefore x(N-x) = kN \quad \therefore N(x-k) = x^2 \text{ 对每个 } i=1, 2, \dots, t$$

由 $p_i | N$, 由 $N(x-k) = x^2 \therefore p_i | x^2 \therefore p_i | x \therefore p_1 p_2 \cdots p_t | x$ 即 $N|x$

$$\therefore x \geq N \quad \therefore y \leq 0 \text{ 矛盾}$$



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5. 设 $n=40|d$ $d \geq 2$ $d \in \mathbb{Z}$. $\therefore$ 对 $d = \frac{n}{40}$. 记 d_1, d_2 (d_1, d_2 为 n 的正因数) 使得 $d = \frac{n}{40} = d_1 \cdot d_2$ ①

设 $n = d_1 u = d_2 v$. $u, v \in \mathbb{Z}_+$ 代入①得 $\frac{n}{40} = \frac{n}{u} - \frac{n}{v}$

$$\Rightarrow \frac{1}{40} = \frac{1}{u} - \frac{1}{v} \Rightarrow u < 40 \mid uv = 40(v-u) \text{ ②}$$

由于 40 是偶数, 由②得 $40 \mid uv \Rightarrow 40 \mid v$. 设 $v = 40k$

$k \in \mathbb{Z}_+$ 代入② $u \cdot 40k = 40(v-u)$

$$\therefore u(k+1) = 40k \text{ ③} \quad \text{由于 } (k, k+1) = 1 \text{ 由③得}$$

$$k+1 \mid 40k \Rightarrow k+1 \mid 40 \quad \therefore k+1 = 40 \quad \therefore k = 39$$

$$\therefore v = 40 \times 39 = 1560 \quad \therefore n \geq v = 1560$$

对 $n = 1560$ 设 n 的所有正因数构成的集合为 $A \cup B$

$$\text{其中 } A = \{1, 2, 4, 5, 8, 10, 16, 20, 25, 40, 50, 80, 100, 200, 400\}$$

$$B = \{a \cdot 40 \mid a \in A\}$$



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(1) $\exists d \in A$ 且 $d \neq 16, 80, 400$, $2 \mid 2d \in A \therefore d = 2d - d$ ✓

$d = 16 \in A$ $2 \mid 16 = 20 - 4$ ~~其中~~ $20, 4 \in A$

$d = 80 \in A$ $2 \mid 80 = 100 - 20$ ~~其中~~ $100, 20 \in A$

$d = 400 \in A$ $2 \mid 400 = 401 - 1$ ~~其中~~ $401 \in B, 1 \in A$

(2) $\forall d \in B$ 设 $d = 401d'$, $d' \in A \setminus \{400\}$ 由 (1) 得

存在 $d_1, d_2 \in A$ 使得 $d' = d_1 - d_2$

$\therefore d = 401d' = 401d_1 - 401d_2$, 其中 $401d_1, 401d_2 \in B$

$\therefore n = 160400$ 是 $\times$ 子数 $\therefore n_{\min} = 160400$

6
$$S_n = \frac{10^n}{n^2 + n^2 + n + 1} = \frac{10^n}{(n+1)(n^2+1)}$$
 $n=1$ 时 $S_1 = \frac{10}{2 \cdot 2} \notin \mathbb{Z}_4 \times$

$n \geq 2$ 时 $\therefore (n+1, n^2+1) = (n+1, 2) = 1$ 或 2

若 n 为偶数, 则 $n+1, n^2+1$ 均为奇数 $\therefore 5 \mid n+1, n^2+1$ 不成立

若 n 为奇数, 若 $5 \nmid n^2+1$, 则 $n^2+1 = 2^l$ ($l \in \mathbb{Z}_+$)



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$$\because n^2+1 \equiv 1+1 \equiv 2 \pmod{4} \quad \therefore l=1 \quad \therefore n=1 \text{ 矛盾}$$

$$\therefore 5 \nmid n^2+1 \quad 2 \parallel n^2+1 \quad \text{设 } n^2+1 = 2 \cdot 5^t, n^2+1 = 2^k \quad (k \geq 2)$$

$t \in \mathbb{Z}_+$ $k \in \mathbb{Z}_+$

$$\text{若 } k=2, n=3 \quad S_3=25 \checkmark$$

$$\text{若 } k \geq 3 \quad (2^{k-1})^2+1 = 2 \cdot 5^t \Rightarrow 2^{2k-2} - 2 \cdot 2^{k-1} + 2 = 2 \cdot 5^t$$

$$\Rightarrow 5^t - 1 = 2^{2k-3} - 2^{k-1} = 2^{k-1}(2^{k-1} - 1) \quad ①$$

5

$$\therefore 5^t \equiv \begin{cases} 1 & t \text{ 为偶数} \\ 5 & t \text{ 为奇数} \end{cases} \pmod{8} \quad \therefore 5^t - 1 \equiv \begin{cases} 0 & t \text{ 为偶数} \\ 4 & t \text{ 为奇数} \end{cases} \pmod{8}$$

$$\text{由于 } 8 \nmid 2^k \quad \therefore \text{由 } ① \text{ 知 } t \text{ 为偶数} \quad \text{设 } t=2m, (m \in \mathbb{Z}_+)$$

$$\text{由 } ① \text{ 得 } (5^m - 1)(5^m + 1) = 2^k (2^{k-1} - 1) \quad ②$$

$$\because 5^m+1 \equiv 1+1 \equiv 2 \pmod{4} \quad \therefore \text{设 } 5^m-1 = 2^{k-1}a, \text{ 其中 } a$$

为奇数代入②得 $2^{k-1}a(2^{k-1}a+2) = 2^k(2^{k-1}-1)$



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$$\Rightarrow a(2^{k-2}a+1)=2^{k-1}-1 \quad (3)$$

若 $a \geq 3$ 则 (3) 左 $> 2^{k-2} \cdot 2 + 1 > (3)$ 右. 矛盾 $\therefore a=1$

$$\text{代入 } (3) \text{ 得 } 2^{k-2} + 1 = 2^{k-1} - 1 \quad 2^{k-2} = 2 \quad \therefore k=3$$

$$\therefore n = 2^3 - 1 = 7. \quad S_7 = 25000 \checkmark$$

9. 设 $\frac{a}{b}, \frac{c}{d}$ 满足 $\frac{a}{b} + \frac{c}{d} = 1$, 且 $a, c, n \in \mathbb{Z}, b, d \in \mathbb{Z}_+$

且 $(a, b) = 1, (c, d) = 1$. 由 $a + \frac{bc}{d} = bn$ 得 $d \mid bc \Rightarrow d \mid b$.

$$(2) \text{ 若 } b \mid d \quad \therefore b = d$$

$$\text{设 } \frac{x^4-1}{y+1} = \frac{a}{b}, a \in \mathbb{Z}, b \in \mathbb{Z}_+, (a, b) = 1, \frac{y^4-1}{x+1} = \frac{c}{d}, c \in \mathbb{Z}, d \in \mathbb{Z}_+, (c, d) = 1$$

$$2. \text{ 若 } b = d \quad \therefore x+1 \mid x^4-1 = (x^2+1)(x+1)(x-1), \quad y+1 \mid y^4-1 = (y^2+1)(y+1)(y-1)$$

$$\therefore \frac{a}{b} \cdot \frac{c}{d} = \frac{x^4-1}{x+1} \cdot \frac{y^4-1}{y+1} \in \mathbb{Z} \quad \therefore b = d = 1$$



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$$\therefore \frac{x^4-1}{y+1} \cdot \frac{y^4-1}{x+1} \in \mathbb{Z} \quad \therefore x+1 \mid y^4-1$$

$$\therefore x^4 y^4 - 1 = x^4 \underbrace{(y^4-1)}_{\downarrow} + \underbrace{x^4-1}$$

$$(y^4-1)(y^{4 \cdot 10} + y^{4 \cdot 9} + y^{4 \cdot 8} + \dots + y^4 + 1)$$

$$\text{即 } y^4-1 \mid y^{44}-1 \quad \therefore x+1 \mid y^{44}-1 \quad \therefore x+1 \mid x^4 y^{44}-1$$

8. 若 $a=1$ 时 对任意 $n \in \mathbb{Z}_+$ 有 $n^2 < n(n+1) < (n+1)^2$ ~~$\therefore n(n+1)$ 不是完全平方数~~
 若 $a=2$ 时 对任意 $n \in \mathbb{Z}_+$ 有 $n^2 < n(n+2) < (n+1)^2$ ~~$\therefore n(n+2)$ 不是完全平方数~~
 若 $a=4$ 时 对任意 $n \in \mathbb{Z}_+$ 有 $n^2 < n(n+4) < (n+2)^2$
 若 $n(n+4)$ 是完全平方数, 则 $n(n+4) = (n+1)^2 \Rightarrow 4n = 2n+1$ 矛盾

$\therefore n(n+4)$ 不是完全平方数 $\checkmark$

若 a 为 2 的幂 设 $a=2^b$ $b \geq 3$ 设 $a=8k$ $k \in \mathbb{Z}_+$ 取 $n=k$
 则 $n(n+a) = n(n+8k) = k(k+8k) = (3k)^2$ 不合题意 $\times$



若 a 不是 2 的幂 设 $a = (2l+1) \cdot k$. $l, k \in \mathbb{Z}_+$. $pn = l^2 \cdot k$

$$n \mid n(n+a) = n(n + (2l+1)k) = l^2 \cdot k (l^2 k + (2l+1)k) = (k l (l+1))^2$$

不合题意. $\times$

9. 设 $S_p = p^2 - 4p + 9$. $p=2$ 时 $S_2 = 9 = 3^2 \checkmark$

$p=3$ 时 $S_3 = 24 \times$

$p \geq 5$ 时 假设 $p^2 - 4p + 9 = n^2$ ($n \in \mathbb{Z}_+$)

$$\therefore (p-2) \mid p(p+2) = (n-3)(n+3) \quad (1)$$

$\therefore p-2 \mid p \cdot (p+2)$ 是积印
 $n \equiv 9$ 奇数 $\therefore$ 一定有一个是 3 的倍数 由 (1) 得 $3 \mid (n-3)(n+3)$

则 $3 \mid n$. 设 $n = 3k$ ($k \in \mathbb{Z}_+$) 代入 (1) 得

$$(p-2) \mid p(p+2) = 9(k-1)(k+1) \quad (2)$$

$$(2) \text{ 得 } p \mid 9(k-1)(k+1) \Rightarrow p \mid (k-1)(k+1) \text{ 要 } p \mid k-1 \text{ 或 } p \mid k+1$$



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(1) 若 $p | k-1$ 设 $k-1 = lp, l \in \mathbb{Z}_+$ 则 $\lambda(2)$ 为

$$(p-2) \cdot p \cdot (p+2) = q \cdot lp (lp+2) \Rightarrow (p-2)(p+2) = ql(lp+2) \quad (2)$$

$$\Rightarrow -4 \equiv 18l \pmod{p} \Rightarrow ql+2 \equiv 0 \pmod{p}$$

$$\therefore ql+2 \geq p \Rightarrow ql \geq p-2 \quad (3), p+2 \geq lp+2 \geq p+2$$

易知 $l=1$. $p = ql+2 = q+2 = 11$. $S_{11} = 36^2 \checkmark$

(2) 若 $p | k+1$ 设 $k+1 = lp, l \in \mathbb{Z}_+$ 则 $\lambda(2)$ 为

$$(p-2)p \cdot (p+2) = q(lp-2) \cdot lp \Rightarrow (p-2)(p+2) = ql(lp-2) \quad (4)$$

$$\Rightarrow -4 \equiv -18l \pmod{p} \Rightarrow ql-2 \equiv 0 \pmod{p}$$

$$\therefore ql-2 \geq p \Rightarrow ql \geq p+2 \quad (5), p-2 \geq lp-2 \geq p-2$$

易知 $l=1$. $p = ql-2 = q-2 = 7$. $S_7 = 18^2 \checkmark$



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